

Mimicking Behaviors in Separated Domains

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Behavior Mimicking

Example

Kid mimicking parent making recipe (unknown to kid):



- Parent acts in real kitchen: real eggs, pans, fire, water,...
- Kid in toy kitchen: plastic eggs, toy pots only (no toy pans), no fire, no water,...
- Kid's goal: make recipe in toy kitchen
- Kid maps real kitchen states/actions to toy kitchen:
 - egg in real pan \Rightarrow toy egg in pot
 - pan on stove \Rightarrow pot on toy stove

Behavior Mimicking

Main features:

- Two agents: A (parent), B (kid)
- A operates in \mathcal{D}_A (real kitchen)
- B operates in \mathcal{D}_B (toy kitchen)
- \mathcal{D}_A and \mathcal{D}_B separated:
 - Actions in \mathcal{D}_A do not affect \mathcal{D}_B , and viceversa
- Mapping φ between behaviors of A and B (on \mathcal{D}_A and \mathcal{D}_B):
 - egg in real pan \Rightarrow toy egg in pot
 - pan on stove \Rightarrow pot on toy stove
- Goal:
 - find a strategy for B to mimic A
 - *Mimicking* defined by φ

Dynamic Domain over propositions $Prop$

- Domain $\mathcal{D} = (S, s_0, \delta, \lambda)$:
 - S finite set of states
 - $s_0 \in S$ initial state
 - $\delta \subseteq S \times S$ transition relation
 - $\lambda : S \mapsto 2^{Prop}$ state-labeling function
- Finite/infinite *traces* as standard: $\tau = s_0 s_1 \dots s_\ell$ (possibly $\ell = \infty$)
- Actions correspond to selecting next transition
- In fact, deterministic

Mimicking Behaviors in Separated Domains (MBSD) Problem Instance

$\mathcal{P} = (\mathcal{D}_A, \mathcal{D}_B, \Phi, Ag_{stop})$, where:

- $\mathcal{D}_A = (S, s_0, \delta^A, \lambda^A)$, dynamic domain over $Prop^A$
- $\mathcal{D}_B = (T, t_0, \delta^B, \lambda^B)$, dynamic domain over $Prop^B$
- $Prop^A \cap Prop^B = \emptyset$
- Φ , mapping specification: LTL_f formula over $Prop^A \cup Prop^B$
- $Ag_{stop} \in \{A, B\}$, designated *stop agent*

Mappings

Intuition: Φ expresses properties of *joint traces* of \mathcal{D}_A and \mathcal{D}_B

$$\tau_A = s_0 s_1 \cdots s_\ell \quad \tau_B = t_0 t_1 \cdots t_\ell \quad \tau_A \cup \tau_B = \begin{pmatrix} s_0 \\ t_0 \end{pmatrix} \begin{pmatrix} s_1 \\ t_1 \end{pmatrix} \cdots \begin{pmatrix} s_\ell \\ t_\ell \end{pmatrix}$$

- Φ : Linear-time temporal formulae over finite traces (LTL_f) over $Prop^A$ and $Prop^B$
- Combines:
 - boolean operators, temporal operators: *next* (**X**), *until* (**U**), always (\square), eventually (\diamond), ...
- can express:
 - $\square\phi$ (always ϕ), $\diamond\phi$ (eventually ϕ), $\square\phi \rightarrow \diamond\psi$ (whenever ϕ eventually ψ), $\phi\mathbf{U}\psi$ (ϕ until ψ), ...

Examples:

- $\Phi = \square(\text{egg_in_real_pan} \rightarrow \text{egg_in_toy_pot}) \wedge \square(\text{pan_on_stove} \rightarrow \text{pot_on_toy_stove})$
- In general, arbitrarily complex mappings: $\square(a \rightarrow \diamond(b\mathbf{U}c)) \wedge \square\diamond q$

Designated Stop Agent

Designated *Stop Agent* decides when to stop

Examples:

- B (parent) announces when recipe is completed
- A (kid) decides when to leave

Choice of Stop Agent strongly affects solution:

- if B (kid) leaves right after starting the game, it trivially mimics B (parent)
- (In kitchen example, parent stops)

Strategies

Strategy for B :

- function $\sigma : S^+ \rightarrow T$
- given sequence of \mathcal{D}_A states, returns move for B (i.e., next \mathcal{D}_B state)

In general, depends on history

As A operates in \mathcal{D}_A and B acts according to a strategy σ , a joint trace is *induced*:

$$\tau_{A,\sigma} = \left(\begin{array}{c} s_0 \\ \sigma(s_0) \end{array} \right) \left(\begin{array}{c} s_1 \\ \sigma(s_0 \ s_1) \end{array} \right) \cdots \left(\begin{array}{c} s_\ell \\ \sigma(s_0 \ s_1 \cdots s_\ell) \end{array} \right)$$

In general, many joint traces exist:

- Result of all choices available to A and consequent B 's responses

Definition (MBSD Solution)

Solution to MBSD problem instance $\mathcal{P} = (\mathcal{D}_A, \mathcal{D}_B, \Phi, Ag_{stop})$:

- Executable strategy σ s.t.:
 - $Ag_{stop} = A$ and for every finite trace τ_A of \mathcal{D}_A , $\tau_{A,\sigma} \models \Phi$; or
 - $Ag_{stop} = B$ and for every infinite trace τ_A of \mathcal{D}_A t.e. finite prefix τ'_A s.t. $\tau'_{A,\sigma} \models \Phi$

Intuition:

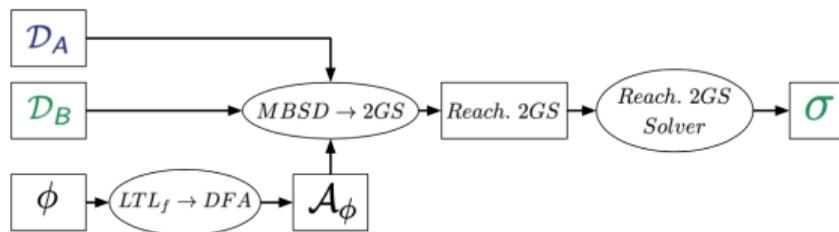
- A Stop Agent: B has a strategy to always keep Φ enforced, no matter how A acts
- B Stop Agent: B has a strategy to enforce Φ at least once, no matter how A acts

Essentially:

- *Synthesis* where environment and system do not affect each other

Solution Approach

- Search for a strategy in a 2-Player game
- Based on constructing DFA \mathcal{A}_ϕ for mapping ϕ (2EXPTIME [De Giacomo&Vardi, 2013])
- DFA \mathcal{A}_ϕ embedded in 2-Player Reachability Game



Theorem

MBSD with general mappings is in:

- *2EXPTIME in combined complexity and mapping complexity*
- *PTIME in domain complexity*

Mapping Classes

Form of mapping affects solution complexity

Two classes of mappings investigated:

- Point-wise Mappings: $\Phi = \bigwedge_{i=1}^k \square(\varphi_i \rightarrow \psi_i)$
 - $\Phi = \square(\text{egg_in_real_pan} \rightarrow \text{egg_in_toy_pot}) \wedge \square(\text{pan_on_stove} \rightarrow \text{pot_on_toy_stove})$
- Target Mappings: $\Phi = \bigwedge_{i=1}^k (\diamond\varphi_i) \rightarrow (\diamond\psi_i)$
 - $\Phi = \diamond\text{salt_added} \rightarrow \diamond\text{talco_added}$
- Recall: properties over \mathcal{D}_A and \mathcal{D}_B are separated:
 - φ_i over Prop^A , ψ_i over Prop^B
 - Prop^A and Prop^B disjoint

Theorem

MBSD with Point-wise mappings is in PTIME for:

- *domain complexity*
- *mapping complexity*
- *combined complexity*

Theorem

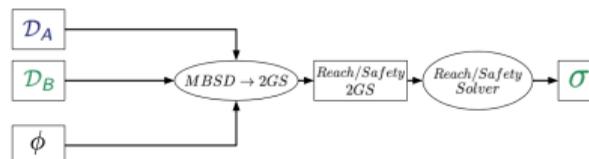
MBSD with Target mappings is

- *PTIME in: domain complexity*
- *PSPACE in:*
 - *mapping complexity*
 - *combined complexity*
- *PSPACE-hard (even with DAG-like \mathcal{D}_A and \mathcal{D}_B)*

Intuition

Each mapping class leads to a different game:

- Point-wise Mappings:
 - Safety Game
 - B 's objective: maintain (continuously) the game in a region where A can be mimicked
 - Game Structure polynomial in size of domains and mapping
- Target Mappings:
 - Reachability Game
 - B 's objective: reach a state where A is (eventually) mimicked
 - Game Structure polynomial in size of domains, exponential in $\#$ of conjuncts in mapping
- Knowing form of mapping saves constructing DFA for ϕ (2EXPTIME)
- Reachability and safety games solvable in PTIME wrt state space of game



Conclusions

Contributions:

- Proposed and formalized MBSD
- General solution approach (2EXPTIME)
- Identified classes of mappings with better computational behavior:
 - Point-wise mappings (PTIME)
 - Target mappings (PSPACE-hard, PTIME wrt domains)
 - Also Tree-like domains (PTIME, not covered in talk)

Open point:

- To what extent separation can yield computational improvements in general, e.g:
- $\Phi = \bigwedge_{i=1}^k \Box(\varphi_i \rightarrow \psi_i)$, with: φ_i LTL_f over $Prop^A$, ψ_i LTL_f over $Prop^B$
- Conjunction of Point-wise and Target-mappings

Thank you!

Questions?